



# Compilers

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## Lexical Specification

- At least one:  $A^+$   $\equiv AA^*$
- Union:  $A \mid B$   $\equiv A + B$
- Option:  $A?$   $\equiv A + \varepsilon$
- Range:  $'a' + 'b' + \dots + 'z'$   $\equiv [a-z]$
- Excluded range:  
complement of  $[a-z]$   $\equiv [^a-z]$

- Last lecture: a specification for the predicate

$$s \in L(R)$$

- Not enough!

## 1. Write a rexp for the lexemes of each token class

- Number =  $\text{digit}^+$
- Keyword = 'if' + 'else' + ...
- Identifier =  $\text{letter}(\text{letter} + \text{digit})^*$
- OpenPar = '('
- ...

2. Construct  $R$ , matching all lexemes for all tokens

$$R = \text{Keyword} + \text{Identifier} + \text{Number} + \dots$$
$$= R_1 + R_2 + \dots$$

3. Let input be  $x_1 \dots x_n$

For  $1 \leq i \leq n$  check

$$x_1 \dots x_i \in L(R)$$

4. If success, then we know that

$$x_1 \dots x_i \in L(R_j) \text{ for some } j$$

5. Remove  $x_1 \dots x_i$  from input and go to (3)

- How much input is used?

- Which token is used?



- What if no rule matches?

- Regular expressions are a concise notation for string patterns
- Use in lexical analysis requires small extensions
  - To resolve ambiguities
  - To handle errors
- Good algorithms known
  - Require only single pass over the input
  - Few operations per character (table lookup)