



Compilers

Static vs. Dynamic Typing

- Static type systems detect common errors
- But some correct programs are disallowed
 - Some argue for dynamic type checking instead
 - Others want more expressive static type checking
- But more expressive type systems are more complex

- The dynamic type of an object is the class **C** that is used in the “**new C**” expression that created it
 - A run-time notion
 - Even languages that are not statically typed have the notion of dynamic type
- The static type of an expression captures all dynamic types the expression could have
 - A compile-time notion

- Soundness theorem: for all expressions E
 $\text{dynamic_type}(E) = \text{static_type}(E)$

In all executions, E evaluates to values of the type inferred by the compiler.

```
class A { ... }  
class B inherits A {...}  
class Main {  
  x:A ← new A;  
  ...  
  x ← new B;  
  ...  
}
```

Choose the static/dynamic type pairs that are correct. For dynamic type, assume execution has halted at line 14.

	<u>Var</u>	<u>Static Type</u>	<u>Dynamic Type</u>
<input type="checkbox"/>	w	Animal	Lion
<input type="checkbox"/>	x	Animal	Pet
<input type="checkbox"/>	y	Pet	Dog
<input type="checkbox"/>	z	Pet	Pet

Static vs. Dynamic

```
1  class Animal { ... }
2  class Pet inherits Animal { ... }
3  class Cat inherits Pet { ... }
4  class Dog inherits Pet { ... }
5  class Lion inherits Animal { ... }
6  class Main {
7      w:Animal <- new Animal;
8      x:Animal <- new Pet;
9      y:Animal <- new Pet;
10     z:Pet <- new Pet;
11     w <- new Lion;
12     y <- new Dog;
13     z <- new Cat;
→ 14     ...
15 }
```

Soundness theorem for the Cool type system:

$$\forall E. \text{dynamic_type}(E) \leq \text{static_type}(E)$$

- All operations that can be used on an object of type C can also be used on an object of type $C' \leq C$
 - Such as fetching the value of an attribute
 - Or invoking a method on the object
- Subclasses only add attributes or methods
- Methods can be redefined but with same type!