



Compilers

Self Type Checking

- `SELF_TYPE`'s meaning depends on the enclosing class

$$O, M, C \vdash e : T$$

An expression e occurring in the body of C has static type T given a variable type environment O and method signatures M

- The next step is to design type rules using **SELF_TYPE**
- Most of the rules remain the same
 - But use the new \leq and **lub**

$$\frac{\begin{array}{l} O(\text{Id}) = T_0 \\ O, M, C \vdash e_1 : T_0 \\ T_1 \leq T_0 \end{array}}{O, M, C \vdash \text{Id} \leftarrow e_1 : T_1}$$

- Recall the old rule for dispatch

$$\begin{array}{c}
 O, M, C \vdash e_0 : T_0 \\
 \vdots \\
 O, M, C \vdash e_n : T_n \\
 M(T_0, f) = (T_1', \dots, T_n', T_{n+1}') \\
 T_{n+1}' \neq \text{SELF_TYPE} \\
 T_i \leq T_i' \quad 1 \leq i \leq n \\
 \hline
 O, M, C \vdash e_0.f(e_1, \dots, e_n) : T_{n+1}'
 \end{array}$$

- If the return type of the method is **SELF_TYPE** then the type of the dispatch is the type of the dispatch expression:

$$\frac{\begin{array}{l} O, M, C \vdash e_0 : T_0 \\ \vdots \\ O, M, C \vdash e_n : T_n \\ M(T_0, f) = (T_1', \dots, T_n', \text{SELF_TYPE}) \\ T_i \leq T_i' \quad 1 \leq i \leq n \end{array}}{O, M, C \vdash e_0.f(e_1, \dots, e_n) : T_0}$$

- Formal parameters cannot be `SELF_TYPE`
- Actual arguments can be `SELF_TYPE`
 - The extended \leq relation handles this case
- The type T_0 of the dispatch expression could be `SELF_TYPE`
 - Which class is used to find the declaration of `f`?
 - Answer: it is safe to use the class where the dispatch appears

- Recall the original rule for static dispatch

$$\begin{array}{l}
 O, M, C \vdash e_0 : T_0 \\
 \vdots \\
 O, M, C \vdash e_n : T_n \\
 T_0 \leq T \\
 M(T, f) = (T_1', \dots, T_n', T_{n+1}') \\
 T_{n+1}' \neq \text{SELF_TYPE} \\
 T_i \leq T_i' \quad 1 \leq i \leq n \\
 \hline
 O, M, C \vdash e_0 @ T.f(e_1, \dots, e_n) : T_{n+1}'
 \end{array}$$

- If the return type of the method is **SELF_TYPE** we have:

$$\begin{array}{l} O, M, C \vdash e_0 : T_0 \\ \vdots \\ O, M, C \vdash e_n : T_n \\ T_0 \leq T \\ M(T, f) = (T_1', \dots, T_n', \text{SELF_TYPE}) \\ \frac{T_i \leq T_i' \quad 1 \leq i \leq n}{O, M, C \vdash e_0 @ T.f(e_1, \dots, e_n) : T_0} \end{array}$$

- Why is this rule correct?
- If we dispatch a method returning `SELF_TYPE` in class `T`, don't we get back a `T`?
- No. `SELF_TYPE` is the type of the self parameter, which may be a subtype of the class in which the method appears

- There are two new rules using **SELF_TYPE**

$$O, M, C \vdash \text{self} : \text{SELF_TYPE}_C$$

$$O, M, C \vdash \text{new SELF_TYPE} : \text{SELF_TYPE}_C$$

Self Type Checking

Choose the static/dynamic type pairs that are correct. For dynamic type, assume execution has halted at line 15.

	<u>Var</u>	<u>Static Type</u>	<u>Dynamic Type</u>
<input type="checkbox"/>	w	Animal	Pet
<input type="checkbox"/>	x	Animal	Lion
<input type="checkbox"/>	y	Pet	Pet
<input type="checkbox"/>	z	Animal	Dog

```
1 class Animal {
2   clone() : SELF_TYPE { new SELF_TYPE }
3 }
4 class Pet inherits Animal {
5   clone() : Pet { new SELF_TYPE }
6 }
7 class Cat inherits Pet { ... }
8 class Dog inherits Pet { ... }
9 class Lion inherits Animal { ... }
10 class Main {
11   w:Animal <- (new Animal).clone();
12   x:Animal <- (new Lion).clone();
13   y:Pet <- (new Cat).clone();
14   z:Animal <- (new Dog)@Animal.clone();
15   ...
16 }
```



- The extended \leq and **lub** operations can do a lot of the work.
- **SELF_TYPE** can be used only in a few places. Be sure it isn't used anywhere else.
- A use of **SELF_TYPE** always refers to any subtype of the current class
 - The exception is the type checking of dispatch. The method return type of **SELF_TYPE** might have nothing to do with the current class

- SELF_TYPE is a research idea
 - It adds more expressiveness to the type system
- SELF_TYPE is itself not so important
 - except for the project
- Rather, SELF_TYPE is meant to illustrate that type checking can be quite subtle
- In practice, there should be a balance between the complexity of the type system and its expressiveness