



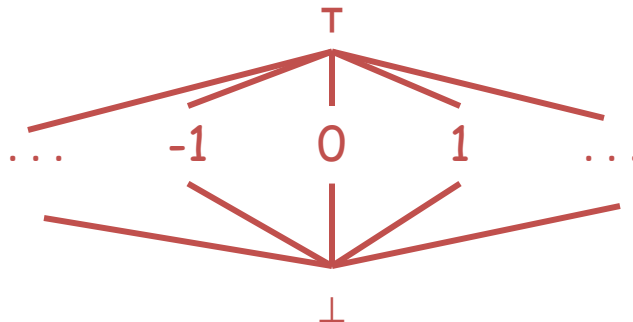
Compilers

Orderings

- We can simplify the presentation of the analysis by ordering the values

$$\perp < c < T$$

- Drawing a picture with “lower” values drawn lower, we get



- \top is the greatest value, \perp is the least
 - All constants are in between and incomparable
- Let *lub* be the least-upper bound in this ordering
- Rules 1-4 can be written using lub:
$$C(s, x, \text{in}) = \text{lub} \{ C(p, x, \text{out}) \mid p \text{ is a predecessor of } s \}$$

- Simply saying “repeat until nothing changes” doesn’t guarantee that eventually nothing changes
- The use of lub explains why the algorithm terminates
 - Values start as \perp and only *increase*
 - \perp can change to a constant, and a constant to \top
 - Thus, $C(s, x, _)$ can change at most twice

Thus the constant propagation algorithm is linear in program size

Number of steps =

Number of $C(\dots)$ values computed $\times 2 =$

Number of program statements $\times 4$